CS152: Programming Languages

Lecture 3 — Operational Semantics

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Where we are

- Done: Caml basics, "IMP" syntax, structural induction
- Now: Operational semantics for our little "IMP" language
  - Most of what you need for Homework 1
  - (But Problem 4 requires proofs over semantics)

Review

IMP's abstract syntax is defined inductively:

$$\begin{array}{rcl} s & ::= & \textbf{skip} \mid x := e \mid s; s \mid \textbf{if}\ e\ s\ s \mid \textbf{while}\ e\ s \\ e & ::= & c \mid x \mid e + e \mid e * e \\ (c & \in & \{\ldots, -2, -1, 0, 1, 2, \ldots\}) \\ (x & \in & \{\mathrm{x}_1, \mathrm{x}_2, \ldots, \mathrm{y}_1, \mathrm{y}_2, \ldots, \mathrm{z}_1, \mathrm{z}_2, \ldots, \ldots\}) \end{array}$$

We haven't yet said what programs *mean*! (Syntax is boring)

Encode our "social understanding" about variables and control flow

Outline

- Semantics for expressions
  1. Informal idea; the need for *heaps*
  2. Definition of heaps
  3. The evaluation *judgment* (a relation form)
  4. The evaluation *inference rules* (the relation definition)
  5. Using inference rules
     - *Derivation trees* as interpreters
     - Or as *proofs* about expressions
  6. *Metatheory*: Proofs about the semantics
- Then semantics for statements
  - ...

Informal idea

Given e, what c does it evaluate to?

$$\mathbf{1} + \mathbf{2} \qquad\qquad \mathtt{x} + \mathbf{2}$$

Informal idea

Given e, what c does it evaluate to?

$$1 + 2 \qquad\qquad \mathtt{x} + 2$$

It depends on the values of variables (of course)

Use a heap H for a total function from variables to constants

- Could use partial functions, but then $\exists$ H and e for which there is no c

We'll define a *relation* over triples of H, e, and c

- Will turn out to be *function* if we view H and e as inputs and c as output
- With our metalanguage, easier to define a relation and then prove it is a function (if, in fact, it is)

Heaps

$$H ::= \cdot \mid H, x \mapsto c$$

A lookup-function for heaps:

$$H(x) = \begin{cases} c & \text{if} \quad H = H', x \mapsto c \\ H'(x) & \text{if} \quad H = H', y \mapsto c' \text{ and } y \neq x \\ 0 & \text{if} \quad H = \cdot \end{cases}$$

- Last case avoids "errors" (makes function *total*)

"What heap to use" will arise in the semantics of statements

- For expression evaluation, "we are given an H"

The judgment

We will write: $\boxed{H \; ; e \Downarrow c}$

to mean, "e evaluates to c under heap H"

It is just a relation on triples of the form (H, e, c)

We just made up metasyntax $H \; ; e \Downarrow c$ to follow PL convention and to distinguish it from other relations

We can write: $., x \mapsto 3 \; ; x + y \Downarrow 3$, which will turn out to be *true*
(this triple will be in the relation we define)

Or: $., x \mapsto 3 \; ; x + y \Downarrow 6$, which will turn out to be *false*
(this triple will not be in the relation we define)

Inference rules

CONST

$$\frac{}{H\ ;\ c \Downarrow c}$$

VAR

$$\frac{}{H\ ;\ x \Downarrow H(x)}$$

ADD

$$\frac{H\ ;\ e_1 \Downarrow c_1 \qquad H\ ;\ e_2 \Downarrow c_2}{H\ ;\ e_1 + e_2 \Downarrow c_1 + c_2}$$

MULT

$$\frac{H\ ;\ e_1 \Downarrow c_1 \qquad H\ ;\ e_2 \Downarrow c_2}{H\ ;\ e_1 * e_2 \Downarrow c_1 * c_2}$$

Top: *hypotheses*
Bottom: *conclusion* (read first)

By definition, if all hypotheses hold, then the conclusion holds

Each rule is a *schema* you "instantiate consistently"

- So rules "work" "for all" H, c, e_1, etc.
- But "each" e_1 has to be the "same" expression

Instantiating rules

Example instantiation:

$$\frac{\cdot, \mathrm{y} \mapsto \mathbf{4} \;;\; \mathbf{3} + \mathrm{y} \Downarrow \mathbf{7} \qquad \cdot, \mathrm{y} \mapsto \mathbf{4} \;;\; \mathbf{5} \Downarrow \mathbf{5}}{\cdot, \mathrm{y} \mapsto \mathbf{4} \;;\; (\mathbf{3} + \mathrm{y}) + \mathbf{5} \Downarrow \mathbf{12}}$$

Instantiates:

$$\text{ADD}\;\frac{\boldsymbol{H} \;;\; \boldsymbol{e_1} \Downarrow \boldsymbol{c_1} \qquad \boldsymbol{H} \;;\; \boldsymbol{e_2} \Downarrow \boldsymbol{c_2}}{\boldsymbol{H} \;;\; \boldsymbol{e_1} + \boldsymbol{e_2} \Downarrow \boldsymbol{c_1} + \boldsymbol{c_2}}$$

with

$\boldsymbol{H} = \cdot, \mathrm{y} \mapsto \mathbf{4}$

$\boldsymbol{e_1} = (\mathbf{3} + \mathrm{y})$

$\boldsymbol{c_1} = \mathbf{7}$

$\boldsymbol{e_2} = \mathbf{5}$

$\boldsymbol{c_2} = \mathbf{5}$

Derivations

A *(complete) derivation* is a tree of instantiations with *axioms* at the leaves

Example:

$$\dfrac{\dfrac{}{\cdot, \mathrm{y}\mapsto\mathbf{4}\ ;\ \mathbf{3} \Downarrow \mathbf{3}} \qquad \dfrac{}{\cdot, \mathrm{y}\mapsto\mathbf{4}\ ;\ \mathrm{y} \Downarrow \mathbf{4}}}{\dfrac{\cdot, \mathrm{y}\mapsto\mathbf{4}\ ;\ \mathbf{3} + \mathrm{y} \Downarrow \mathbf{7}}{}} \qquad \dfrac{}{\cdot, \mathrm{y}\mapsto\mathbf{4}\ ;\ \mathbf{5} \Downarrow \mathbf{5}}$$

$$\overline{\cdot, \mathrm{y}\mapsto\mathbf{4}\ ;\ (\mathbf{3} + \mathrm{y}) + \mathbf{5} \Downarrow \mathbf{12}}$$

By definition, $\boldsymbol{H}\ ;\ \boldsymbol{e} \Downarrow \boldsymbol{c}$ if there exists a derivation with $\boldsymbol{H}\ ;\ \boldsymbol{e} \Downarrow \boldsymbol{c}$ at the root

Back to relations

So what relation do our inference rules define?

- Start with empty relation (no triples) R_0
- Let R_i be R_{i-1} union all $H \,;\, e \Downarrow c$ such that we can instantiate some inference rule to have conclusion $H \,;\, e \Downarrow c$ and all hypotheses in R_{i-1}
  - So R_i is all triples at the bottom of height-j complete derivations for $j \leq i$
- R_∞ is the relation we defined
  - All triples at the bottom of complete derivations

For the math folks: R_∞ is the smallest relation closed under the inference rules

What are these things?

We can view the inference rules as defining an *interpreter*

- Complete derivation shows recursive calls to the "evaluate expression" function
  - Recursive calls from conclusion to hypotheses
  - *Syntax-directed* means the interpreter need not "search"

- See OCaml code in Homework 1

Or we can view the inference rules as defining a *proof system*

- Complete derivation proves facts from other facts starting with axioms
  - Facts established from hypotheses to conclusions

Some theorems

- Progress: For all H and e, there exists a c such that $H\ ;\ e \Downarrow c$.
- Determinacy: For all H and e, there is at most one c such that $H\ ;\ e \Downarrow c$.

We rigged it that way...
what would division, undefined-variables, or `gettime()` do?

Proofs are by induction on the the structure (i.e., height) of the expression e.

On to statements

A statement doesn't produce a constant.

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We could define $H_1 \; ; s \Downarrow H_2$

- Would be a partial function from H_1 and s to H_2
- Works fine; could be a homework problem

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A statement doesn't produce a constant.

It produces a new, possibly-different heap.

- If it terminates

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- Would be a partial function from H_1 and s to H_2
- Works fine; could be a homework problem

Instead we'll define a "small-step" semantics and then "iterate" to "run the program"

$$\boxed{H_1 \; ; s_1 \rightarrow H_2 \; ; s_2}$$

Statement semantics

$$\boxed{H_1 \; ; s_1 \rightarrow H_2 \; ; s_2}$$

ASSIGN
$$\frac{H \; ; e \Downarrow c}{H \; ; x := e \rightarrow H, x \mapsto c \; ; \textsf{skip}}$$

SEQ1
$$\frac{}{H \; ; \textsf{skip}; s \rightarrow H \; ; s}$$

SEQ2
$$\frac{H \; ; s_1 \rightarrow H' \; ; s_1'}{H \; ; s_1; s_2 \rightarrow H' \; ; s_1'; s_2}$$

IF1
$$\frac{H \; ; e \Downarrow c \qquad c > 0}{H \; ; \textsf{if } e \; s_1 \; s_2 \rightarrow H \; ; s_1}$$

IF2
$$\frac{H \; ; e \Downarrow c \qquad c \leq 0}{H \; ; \textsf{if } e \; s_1 \; s_2 \rightarrow H \; ; s_2}$$

Statement semantics cont'd

What about **while** e s (do s and loop if $e > 0$)?

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$$\frac{\textsc{while}}{H\ ;\ \textbf{while}\ e\ s \rightarrow H\ ;\ \textbf{if}\ e\ (s; \textbf{while}\ e\ s)\ \textbf{skip}}$$

Many other equivalent definitions possible

Program semantics

Defined $H \; ; \; s \rightarrow H' \; ; \; s'$, but what does "$s$" mean/do?

Our machine iterates: $H_1;s_1 \rightarrow H_2;s_2 \rightarrow H_3;s_3 \ldots$,
with each step justified by a complete derivation using our single-step statement semantics

Let $H_1 \; ; \; s_1 \; \rightarrow^n \; H_2 \; ; \; s_2$ mean "becomes after n steps"

Let $H_1 \; ; \; s_1 \rightarrow^* H_2 \; ; \; s_2$ mean "becomes after 0 or more steps"

Pick a special "answer" variable `ans`

The program s produces c if $\cdot \; ; \; s \rightarrow^* H \; ; \;$ **skip** and $H(\texttt{ans}) = c$

Does every s produce a c?

Example program execution

$\mathrm{x} := \mathbf{3}; (\mathrm{y} := \mathbf{1}; \textbf{while } \mathrm{x}\ (\mathrm{y} := \mathrm{y} * \mathrm{x}; \mathrm{x} := \mathrm{x} - \mathbf{1}))$

Let's write some of the state sequence. You can justify each step with a full derivation. Let $s = (\mathrm{y} := \mathrm{y} * \mathrm{x}; \mathrm{x} := \mathrm{x} - \mathbf{1})$.

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$$\rightarrow \quad \cdot, \mathrm{x} \mapsto \mathbf{3}; \textbf{skip}; \mathrm{y} := \mathbf{1}; \textbf{while}\ \mathrm{x}\ s$$

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$$\begin{array}{ll} & \cdot; \mathrm{x} := 3; \mathrm{y} := 1; \textbf{while } \mathrm{x}\ s \\ \rightarrow & \cdot, \mathrm{x} \mapsto 3; \textbf{skip}; \mathrm{y} := 1; \textbf{while } \mathrm{x}\ s \\ \rightarrow & \cdot, \mathrm{x} \mapsto 3; \mathrm{y} := 1; \textbf{while } \mathrm{x}\ s \end{array}$$

Example program execution

$$x := 3; (y := 1; \textbf{while } x\ (y := y * x; x := x - 1))$$

Let's write some of the state sequence. You can justify each step with a full derivation. Let $s = (y := y * x; x := x - 1)$.

$$\cdot; x := 3; y := 1; \textbf{while } x\ s$$

$$\rightarrow \quad \cdot, x \mapsto 3; \textbf{skip}; y := 1; \textbf{while } x\ s$$

$$\rightarrow \quad \cdot, x \mapsto 3; y := 1; \textbf{while } x\ s$$

$$\rightarrow^2 \quad \cdot, x \mapsto 3, y \mapsto 1; \textbf{while } x\ s$$

Example program execution

$$\text{x} := 3; (\text{y} := 1; \textbf{while } \text{x } (\text{y} := \text{y} * \text{x}; \text{x} := \text{x}-1))$$

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$$\rightarrow \quad \cdot, \text{x} \mapsto 3; \textbf{skip}; \text{y} := 1; \textbf{while } \text{x } s$$

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$$\rightarrow \quad \cdot, \text{x} \mapsto 3, \text{y} \mapsto 1; \textbf{if } \text{x } (s; \textbf{while } \text{x } s) \textbf{ skip}$$

Example program execution

$\mathrm{x} := 3; (\mathrm{y} := 1; \textsf{while}\ \mathrm{x}\ (\mathrm{y} := \mathrm{y} * \mathrm{x}; \mathrm{x} := \mathrm{x}{-}1))$

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$$
\begin{array}{ll}
 & \cdot; \mathrm{x} := 3; \mathrm{y} := 1; \textsf{while}\ \mathrm{x}\ s \\
\rightarrow & \cdot, \mathrm{x} \mapsto 3; \textsf{skip}; \mathrm{y} := 1; \textsf{while}\ \mathrm{x}\ s \\
\rightarrow & \cdot, \mathrm{x} \mapsto 3; \mathrm{y} := 1; \textsf{while}\ \mathrm{x}\ s \\
\rightarrow^2 & \cdot, \mathrm{x} \mapsto 3, \mathrm{y} \mapsto 1; \textsf{while}\ \mathrm{x}\ s \\
\rightarrow & \cdot, \mathrm{x} \mapsto 3, \mathrm{y} \mapsto 1; \textsf{if}\ \mathrm{x}\ (s; \textsf{while}\ \mathrm{x}\ s)\ \textsf{skip} \\
\rightarrow & \cdot, \mathrm{x} \mapsto 3, \mathrm{y} \mapsto 1; \mathrm{y} := \mathrm{y} * \mathrm{x}; \mathrm{x} := \mathrm{x} - 1; \textsf{while}\ \mathrm{x}\ s
\end{array}
$$

Continued...

$$\rightarrow^{2} \quad \cdot, \mathrm{x} \mapsto \mathbf{3}, \mathrm{y} \mapsto \mathbf{1}, \mathrm{y} \mapsto \mathbf{3};\ \mathrm{x} := \mathrm{x} - \mathbf{1};\ \textsf{\textbf{while}}\ \mathrm{x}\ s$$

Continued...

$$\rightarrow^{2} \quad \cdot, \mathrm{x} \mapsto \mathbf{3}, \mathrm{y} \mapsto \mathbf{1}, \mathrm{y} \mapsto \mathbf{3}; \mathrm{x} := \mathrm{x}{-}\mathbf{1}; \textsf{\textbf{while}}\ \mathrm{x}\ s$$

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Continued...

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$$\rightarrow \quad \ldots, \mathrm{y} \mapsto \mathbf{3}, \mathrm{x} \mapsto \mathbf{2}; \textbf{ if } \mathrm{x}\ (s; \textbf{while } \mathrm{x}\ s) \textbf{ skip}$$

Continued...

$$\rightarrow^{2} \quad \cdot, \mathrm{x} \mapsto \mathbf{3}, \mathrm{y} \mapsto \mathbf{1}, \mathrm{y} \mapsto \mathbf{3}; \ \mathrm{x} := \mathrm{x} - \mathbf{1}; \textsf{\textbf{while}}\ \mathrm{x}\ s$$

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$\ldots$

Continued...

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$\rightarrow \quad \ldots, \mathrm{y} \mapsto \mathbf{3}, \mathrm{x} \mapsto \mathbf{2}; \textsf{if}\ \mathrm{x}\ (s; \textsf{while}\ \mathrm{x}\ s)\ \textsf{skip}$

$\ldots$

$\rightarrow \quad \ldots, \mathrm{y} \mapsto \mathbf{6}, \mathrm{x} \mapsto \mathbf{0}; \textsf{skip}$

Where we are

Defined $H \; ; e \Downarrow c$ and $H \; ; s \rightarrow H' \; ; s'$ and extended the latter to give s a meaning

- The way we did expressions is "large-step operational semantics"
- The way we did statements is "small-step operational semantics"
- So now you have seen both

Definition by interpretation: program means what an interpreter (written in a metalanguage) says it means

- Interpreter represents a (very) abstract machine that runs code

Large-step does not distinguish errors and divergence

- But we defined IMP to have no errors
- And expressions never diverge

Establishing Properties

We can prove a property of a terminating program by "running" it.

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$\cdot\ ;\ s\ \rightarrow^{n}\ H\ ;\ $**skip** cannot be derived.

Example: **while** 1 **skip**

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We can prove a property of a terminating program by "running" it.

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Example: **while** $\mathbf{1}$ **skip**

By induction on n, but requires a *stronger induction hypothesis*.

More General Proofs

We can prove properties of executing all programs (satisfying another property)

Example: If H and s have no negative constants and $H\ ;\ s \rightarrow^* H'\ ;\ s'$, then H' and s' have no negative constants.

Example: If for all H, we know s_1 and s_2 terminate, then for all H, we know $H;(s_1; s_2)$ terminates.