

Introduction

CS 152 (Spring 2024)

Harvard University

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Programming Languages

- ▶ More than a catalog of languages and what they can be used for.
- ▶ In this class: foundations of programming languages, the underlying concepts and principles that go into designing and implementing programming languages.
- ▶ How can you learn new languages? How can you design effective languages?

Why?

- ▶ give you the concepts to more easily learn new languages
- ▶ ... and to design and implement new languages
- ▶ golden age of PL
- ▶ elegant math

Aspects of a Language

syntax the structure of its programs

semantics the meaning of its programs

A formal semantics...

- ▶ can be simpler than an implementation, more precise than intuition
- ▶ can answer: is this implementation correct
- ▶ supports the definition of analyses and transformations
 - ▶ prove properties about the language
 - ▶ prove properties about programs written in the language
- ▶ promotes better language design
 - ▶ better understand impact on design decisions
 - ▶ apply semantic insights to improve language

Cool: Type safety

Example: Rust

Rust is memory safe (no dereferencing of null pointers, no dangling pointers), but performance is comparable to C and C++. Lots of memory checking is done statically. Achieves this using a sophisticated type system, with parametric polymorphism and linear types. All at compile time, with no run-time overhead.

Cool: Certified compilers

- ▶ Formal proof that the native code output by CompCert has the same semantics as the original C program.
- ▶ Researchers found zero bugs in the verified part of CompCert vs hundreds of bugs in LLVM and GCC.

Cool: Program Synthesis

Cool: Program Verification

Cool: Differentiable Programming

Cool: Probabilistic Programming

ToC

- ▶ semantics
- ▶ lambda calculus
- ▶ types
- ▶ reasoning about programs
- ▶ misc. topics

Semantics of Programming Languages

Give *mathematical meaning* to programs.

Why *mathematical*?

- ▶ Less ambiguous.
- ▶ More concise.
- ▶ Formal arguments.

Semantics

Styles of Semantics

Operational Semantics

Denotational Semantics

Axiomatic Semantics

Algebraic Semantics

Operational Semantics

Small-Step

Large-Step

Small-Step Operational Semantics

step from configuration to configuration:

$$c_0 \longrightarrow c_1 \longrightarrow \dots \longrightarrow c_n$$

Large-Step Operational Semantics

one step from initial configuration to final answer:

$$c \Downarrow a$$

Denotational Semantics

interpret in mathematical domain

$$[[\text{term}]] = \text{number}$$

$$[[e_1 + e_2]] = [[e_1]] + [[e_2]]$$

...

Axiomatic Semantics

$$\{Pre\} c \{Post\}$$

Algebraic Semantics

Abstract Syntax

Abstract Syntax

$x, y, z \in \mathbf{Var}$

$n, m \in \mathbf{Int}$

$e \in \mathbf{Exp}$

Abstract Syntax

$$x, y, z \in \mathbf{Var}$$

Var is the set of program variables (e.g., foo, bar, baz, i, etc.).

Abstract Syntax

$n, m \in \mathbf{Int}$

Int is the set of constant integers (e.g., 42, -40, 7).

Abstract Syntax

$e \in \mathbf{Exp}$

Exp is the domain of expressions, which we specify using a BNF (Backus-Naur Form) grammar.

Simple Expressions

$$e ::= x$$
$$| n$$
$$| e_1 + e_2$$
$$| e_1 \times e_2$$

Abstract Syntax Tree

$1 + 2 \times 3$

Abstract Syntax Tree

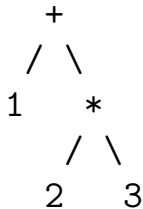
$$1 + 2 \times 3$$

$$1 + (2 \times 3)$$

$$(1 + 2) \times 3$$

Abstract Syntax Tree

$1 + 2 \times 3$

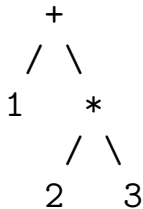


$1 + (2 \times 3)$

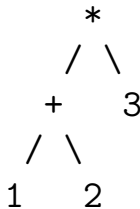
$(1 + 2) \times 3$

Abstract Syntax Tree

$1 + 2 \times 3$



$1 + (2 \times 3)$



$(1 + 2) \times 3$

Def. and Use of Abstract Syntax

- ▶ in OCaml
- ▶ in Coq
- ▶ in Dafny